

Overview: Proof techniques

Direct proof

You go in logical steps from the assumption to the conclusion.

Counter-examples

To disprove a "for all" statement, you provide a single instance for which the property is not true.

Proof by contradiction

You assume that the negation is true, and come to a contradiction.

Contrapositive proof

Instead of $p \rightarrow q$, you prove $(\text{not } q) \rightarrow (\text{not } p)$

Biconditional proof

To prove a $p \leftrightarrow q$ statement, you need to prove $p \rightarrow q$ AND $q \rightarrow p$

TODAY: Proof by induction

Book: Chapter 1, section 1.6

Q: Can you calculate the sum $1+2+3+\dots+(n-1)+n$
(with n a natural number)?

• $sum = 0$

for ($i=1, i++, i=n$)
($sum = sum+i$)

•
$$\sum_{i=1}^n i = 1 + 2 + \dots + n$$
$$= \frac{n}{2} (n+1)$$

Mathematical induction

• proof technique for statements of the form $(\forall n \geq N, n \in \mathbb{Z}) P(n)$

Two steps:

1. Base case: proof of $P(N)$

(the statement holds for the first number, the first domino falls)

2. Induction step: proof of (for all $n \geq N$) $(P(n) \rightarrow P(n+1))$

(If the n -th domino falls, so does the $(n+1)$ -th)

• $P(N)$ is true : Base case

• $P(n) \rightarrow P(n+1) \forall n$ (induction step)

$$\underline{P(N), P(N) \rightarrow P(N+1)} \quad \Rightarrow P(N+1)$$

$$P(2) \rightarrow P(3)$$

$$P(3) \rightarrow P(4)$$

⋮

$$(\forall n \in \mathbb{N}) \left(\sum_{k=1}^n k = \frac{n(n+1)}{2} \right)$$

• Base case . Proof of $P(1)$ statement is true for $n=1$
 $\sum_{k=1}^1 = 1 = \frac{1(1+1)}{2} = 1 \quad \checkmark$

• Induction step . Proof of $\forall m \geq 1, P(m) \rightarrow P(m+1)$

$$P(m) : \sum_{k=1}^m k = \frac{m(m+1)}{2} \quad (\text{assumption})$$

$$P(m+1) : \sum_{k=1}^{m+1} k = \frac{(m+1)(m+2)}{2} \quad (\text{to prove})$$

$$\begin{aligned} \sum_{k=1}^{m+1} k &= (1+2+\dots+m) + m+1 = \frac{m(m+1)}{2} + (m+1) \\ &= (m+1) \left(\frac{m}{2} + 1 \right) \\ &= (m+1) \frac{(m+2)}{2} \quad \square \end{aligned}$$

$\forall n \in \mathbb{N} : 4^n - 1$ is divisible by 3

Proof by induction

• Base case : $P(1)$ is true. $4^1 - 1 = 3$ 3 is divisible by 3 ✓

• Induction step $P(x) \rightarrow P(x+1) \quad \forall x \in \mathbb{N}$

$P(x) : 4^x - 1$ is divisible by 3 (induction hypothesis)

$P(x+1) : 4^{x+1} - 1$ is divisible by 3

$$\begin{aligned} 4^{x+1} - 1 &= 4 \cdot 4^x - 1 = 3 \cdot 4^x + \underbrace{4^x - 1}_{3k \text{ (assumption)}} \\ &= 3(4^x + k) \rightarrow \text{multiple of 3} \quad \checkmark \end{aligned}$$

$$4 \cdot 4^x = (3+1) \cdot 4^x = 3 \cdot 4^x + 4^x$$

A country has two types of coins: 3 cent and 7 cent. Prove that you can make all sums of money starting from 12 cent (up to 1 cent accuracy) from these coins.

$$\begin{aligned}
 12 &= 3 + 3 + \textcircled{3 + 3} \\
 13 &= \textcircled{3 + 3} + \textcircled{7} \\
 14 &= \underline{7 + 7} \\
 15 &= 3 + 3 + 3 + \textcircled{3 + 3} \\
 16 &= \textcircled{3 + 3} + 3 + \textcircled{7} \\
 17 &= \textcircled{7 + 7} + 3 \\
 18 &= 6 \cdot 3 \\
 19 &= 4 \cdot 3 + 7 \\
 &\vdots
 \end{aligned}$$

$$\begin{aligned}
 &3 + 3 \\
 &\leftarrow + \\
 &7
 \end{aligned}$$

$$2 \cdot 7 \rightarrow 5 \cdot 3$$

- 1) if we have 2×3 (at least), we can go 1c up
 - 2) if we have 2×7 , (at least), we can go 1c up
 - 3) we always have 2×7 or 2×3
- (negation: max 1×7 , max 1×3 . cannot add up to 12 or more)

Proof by induction.

- Base case. $P(12)$ holds.

$$12 = 3 + 3 + 3 + 3 \quad \checkmark$$

- Induction step : $P(q) \rightarrow P(q+1)$

1) q has $\underbrace{2 \text{ coins of } 7}_r$ or $\underbrace{2 \text{ coins of } 3}_{s} \text{ at least}$

$\hookrightarrow \neg(r \vee s) = \neg r \wedge \neg s$ is impossible, as $q \geq 12$

2) if r is true, then we can replace 2×7 by 5×3

3) if s is true, then we can replace 2×3 by 7

\hookrightarrow in both cases, we can make $q+1$

$$\forall x \in \mathbb{N}, \exists y \in \mathbb{N} : x^y < y^x$$

negation.

$$\exists x \forall y : x^y \geq y^x$$

→ this statement is false, so we prove the negation.

→ how to start?

try $x=1$. Then $1^2 < 2^1$ → does not work.

$x=2$. $2^3 < 3^2$ → does not work.

$x=3$... $2^3 < 3^2$, $3^3 \leq 3^3$, $4^3 < 3^4$, $5^3 < 3^5$

↳ this might work!

→ so, we try to prove : $\forall y \in \mathbb{N} : y^3 \leq 3^y$

→ we already know : $1^3 \leq 3^1$ (ok for $y=1$)

$2^3 \leq 3^2$ (ok for $y=2$)

$3^3 \leq 3^3$ (ok for $y=3$)

↳ we use induction for $y \geq 3$

• Base case: $y = 3$. $3^3 \leq 3^3$ is true ✓

• Induction step: $P(y) \rightarrow P(y+1) \quad \forall y \geq 3$

$P(y)$: $y^3 \leq 3^y$ — induction hypothesis

$P(y+1)$: $(y+1)^3 \leq 3^{y+1}$
LHS RHS

strategy: we go from LHS to RHS using the induction hypothesis.

$$(y+1)^3 = y^3 + 3y^2 + 3y + 1$$

$$\leq 3^y + 3y^2 + 3y + 1$$

$$\leq 3^y + y^3 + 3y + 1$$

$$\leq 3^y + 3^y + 3y + 1$$

$$\leq 3^y + 3^y + y^3$$

$$\leq 3^y + 3^y + 3^y$$

$$\leq 3 \cdot 3^y$$

$$\leq 3^{y+1}$$



(induction hypothesis)
(since $y \geq 3$, $y^3 \gg 3y^2$)

(induction hyp.)

$$(y^3 \geq 3y + 1 \Leftrightarrow y^3 - 3y \geq 1)$$

$$\Leftrightarrow y(y^2 - 3) \geq 1$$

since $y \geq 3$

End of chapter 1:
Logic and proofs.

We already completed $1/3$ of the
course.

Tomorrow: set theory